

A Semantics for Means-End Relations

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Outline

- 1 Means-end relations in practical reasoning
 - von Wright's example
 - Initial informal analysis

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 - Introduction to PDL
 - Sufficient means-end relations

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 - von Wright's necessary means
 - Involvement

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The conclusion is an action or an intention.

von Wright's Practical Inference

A working example from von Wright.



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- Expression of an agent's desire,
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- Concludes in a necessary action.

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Features:

- Conclusion is necessary on pain of practical irrationality.

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How do the premises make the conclusion necessary?

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We focus on the semantics of means-end relations.

Initial steps for a means-end semantics

- An end is some desirable condition – a proposition.

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- **Objects as means?**

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PDL syntax

Propositional Dynamic Logic is a logic of actions.

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Intuition for $[\alpha]\varphi$: After doing α , φ will hold.

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$w \models [\alpha]\varphi$ iff whenever $w \xrightarrow{\alpha} w'$, then $w' \models \varphi$.

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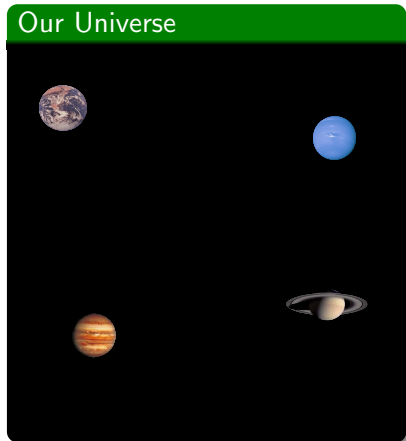
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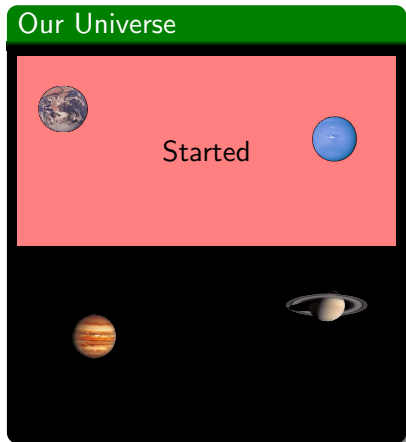
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A simple example of possible worlds



A set of worlds involving
a footrace and starter pistol.

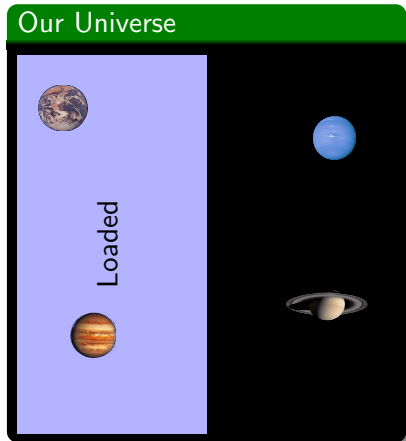
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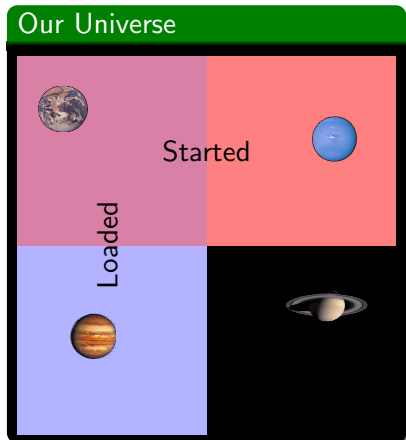
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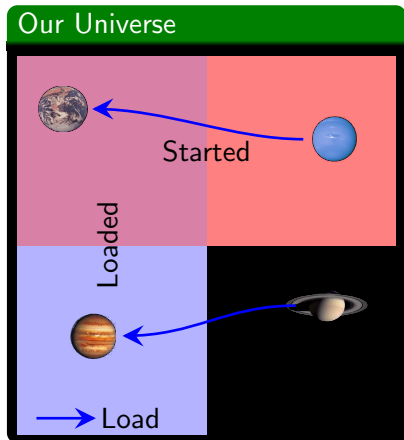
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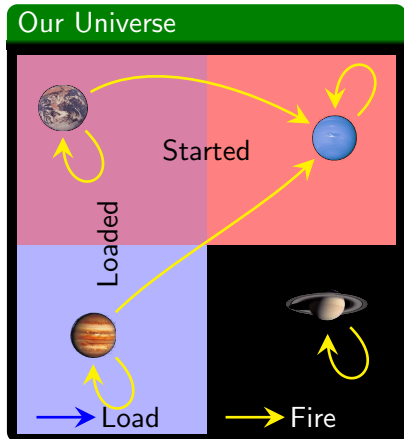
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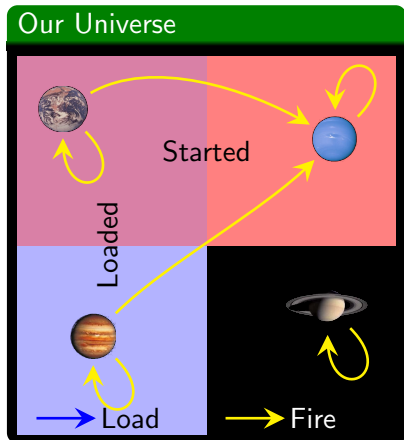
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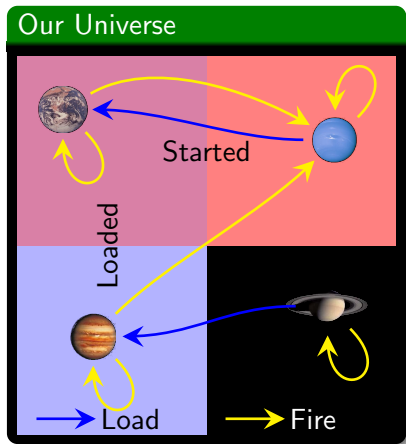
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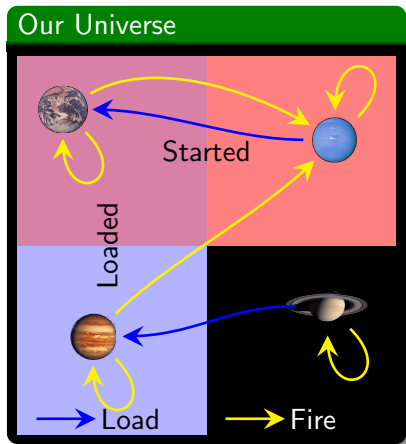
Note: “Fire” means “pull trigger”. We allow misfires.

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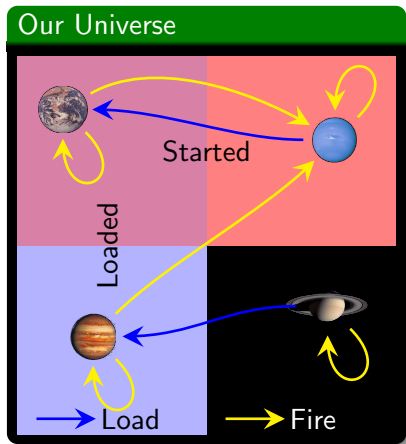
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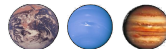
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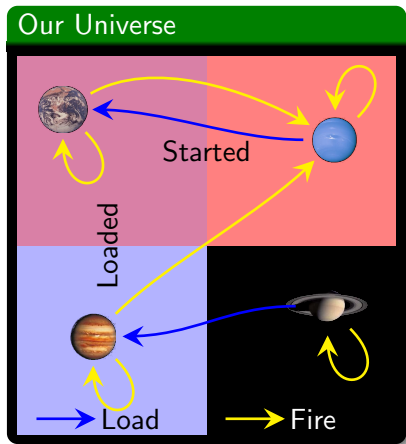
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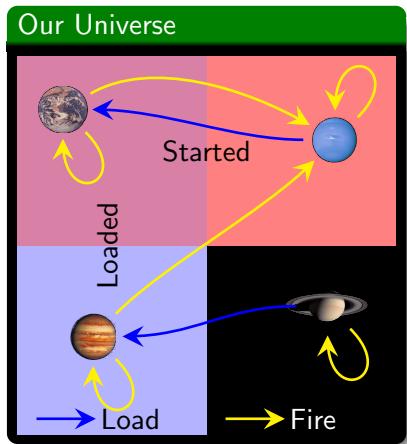
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Sufficient means-end relations

In w , α is a strongly sufficient means to φ

Doing α in w will yield φ

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Doing α in w will yield φ \updownarrow and one can do α in w .

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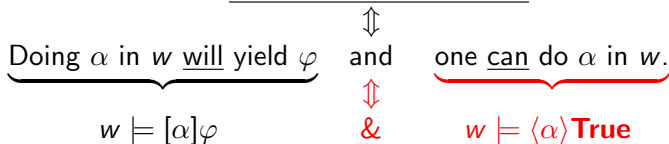
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If α is atomic, then β involves α iff α is a subterm of β .

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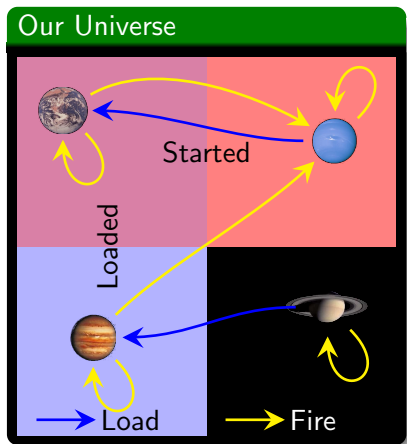
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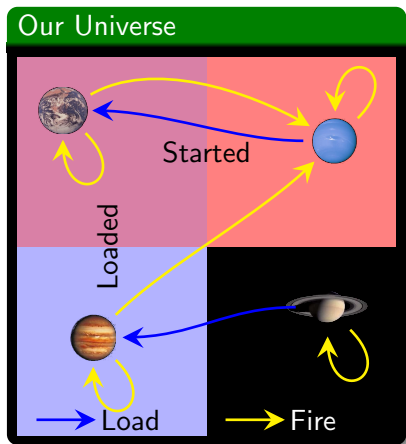
But what if $\alpha = \alpha_1; \alpha_2$?


A necessary means-end relation



In  , load; fire is a necessary means to **Started**.

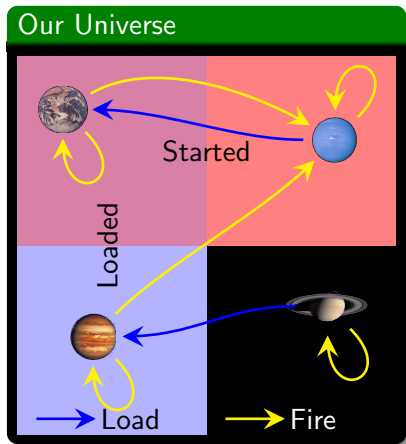
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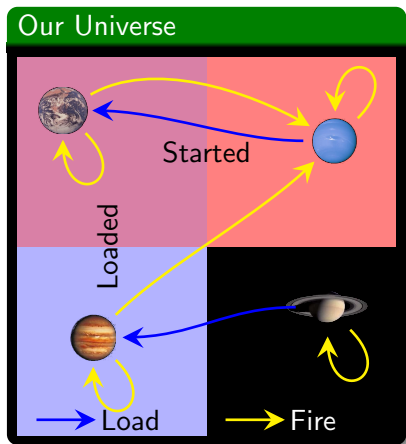
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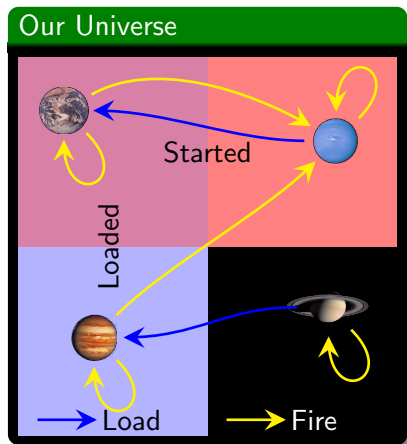



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Does **fire; load; fire; fire** involve **load; fire**? Yes!

Does **skip; load; skip; fire** involve **load; fire**?

A necessary means-end relation



In  , load; fire is a necessary means to **Started**.

Does fire; load; fire; fire involve load; fire? Yes!

Does skip; load; skip; fire involve load; fire? Also yes!

Involvement

Write $\beta \preccurlyeq \alpha$ for: β involves α .

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Basic properties:

- \preceq is a pre-order.
- Non-deterministic choice \cup is the join for \preceq .
- If $\beta \preceq \alpha$, then $\beta; \gamma \preceq \alpha; \gamma$ and $\gamma; \beta \preceq \gamma; \alpha$.

Doing α and involvement

Necessary means:

If α is a necessary means to φ , then

- φ will not be realized without doing α and
- there is a means to realize φ that involves doing α .

α is a necessary means to φ in w iff

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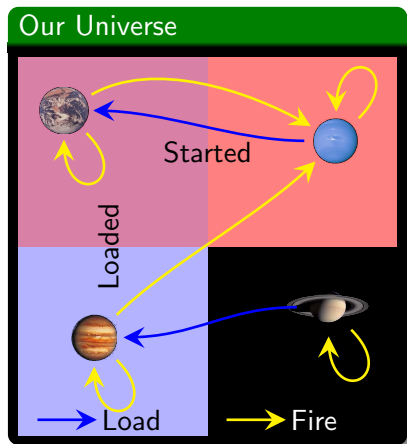
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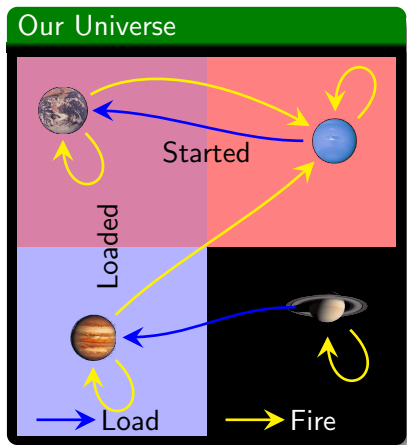
- if $w \models \langle \beta \rangle \varphi$ then $\beta \preceq \alpha$;
- there is a $\beta \preceq \alpha$ such that $w \models \langle \beta \rangle \varphi$.




Examples of necessary means to **Started**



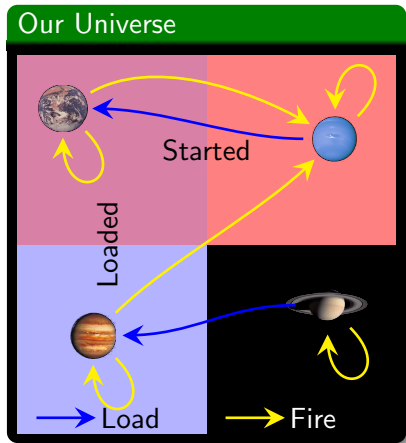
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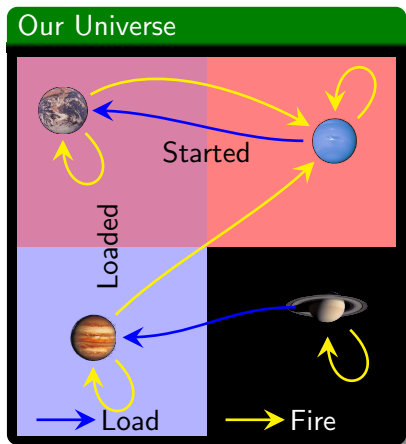
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Examples of necessary means to **Started**



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To realize **Started**,
 one must do some β involving
 every necessary means.

Additional topics

- Involvement with test actions (in paper).

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- **Defined negations for actions.**

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- From means-end relations to artifactual functions.

Outline

- 4 Conditionals and the frame problem
 - Local vs. conditional relations
 - Non-monotonicity

Local means-end relations

Our definition

In w , α is a sufficient means to φ iff
 $w \models [\alpha]\varphi$ & $\langle \alpha \rangle \mathbf{True}$.

This is a very narrow sense of means-end relation.

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Example

“Riding the train is a means to reaching Delft.”



Local means-end relations

Our definition

In w , α is a sufficient means to φ iff
 $w \models [\alpha]\varphi$ & $\langle \alpha \rangle \mathbf{True}$.

This is a very narrow sense of means-end relation.

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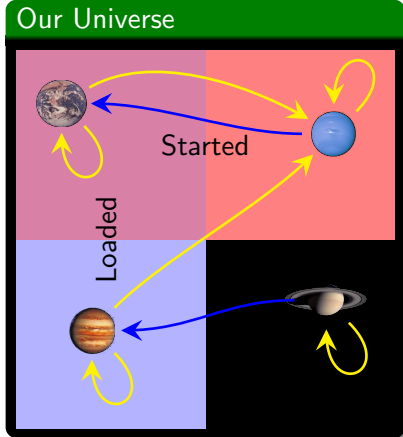
Do we mean this is true just in

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- every world?
- every world in which we are in Eindhoven?
- every “normal” world in which we are in Eindhoven?

Introducing conditional means-end relations

Conditional relation:

Assuming ψ ,
 α is a means to φ .

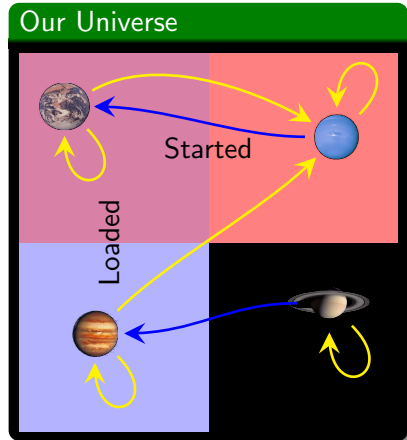


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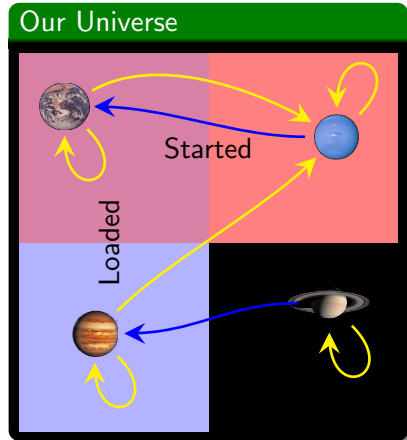
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$\models \psi \rightarrow [\alpha]\varphi$ iff $w \not\models \psi$ or
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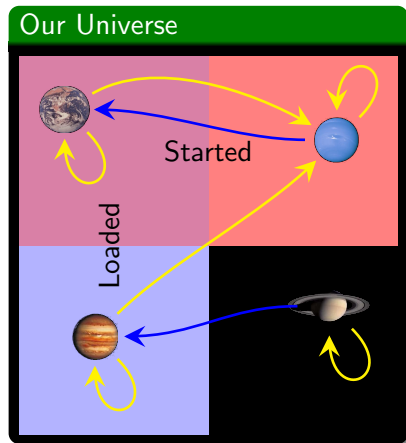
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In every world satisfying ψ ,
 α is a local means to φ .

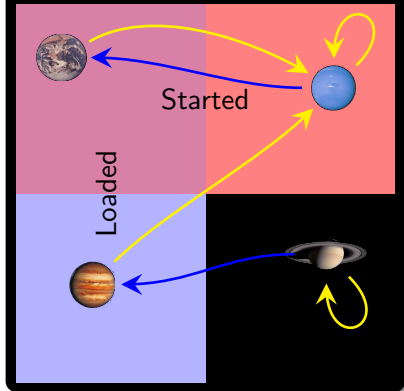


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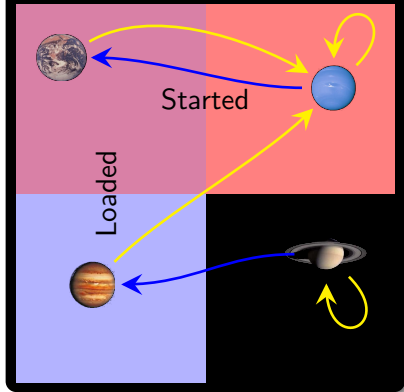


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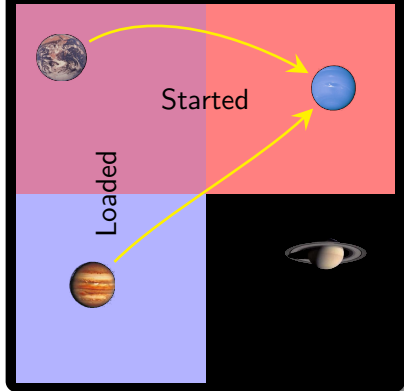


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Reevaluating material implication

(or “Why means-end reasoning is hard”)

A simple derivation:

If I had money, she would marry me.



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Problem: If I rob her, she will hate me and
(**money** & **HATE**) ↗ [propose]**marry**.

Our conditional should be non-monotonic

Non-monotonicity

money \rightarrow [propose]marry but
(money & **HATE**) \nrightarrow [propose]marry.

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Reasoning about means is hard.