Means-end Relations and a Measure of Efficacy

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Outline

Means-end relations

- Interest I: Practical syllogisms
- Interest II: Functional ascriptions
- Propositional Dynamic Logic

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2 Efficacy via fuzzy logic

- Reliability as a fuzzy operator
- The resulting fuzzy logic

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The conclusion is an *action* or an *intention*.

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The conclusion is an *action* or an *intention*.

This premise is a means-end relation.

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Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

An example from von Wright



I want to make the hut habitable.

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I want to make the hut habitable. Unless I heat the hut, it will not be habitable.

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Therefore I must heat the hut.

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• Expression of an agent's desire,

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Note: distinct premises

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- Expression of an agent's desire, Note: distinct premises
- A necessary means-end relation,⁴
- Concludes in a *necessary* action.

But necessary means-end relations are a bit tricky.

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An example from von Wright



I want to make the hut habitable. If I heat the hut, it will be habitable. Therefore, I have reason to heat the hut.

An alternative with a sufficient means-end relation.



• "The function of the heart is to pump blood."

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- "The function of the heart is to pump blood."
- "That switch mutes the television."

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We ascribe functions to biological stuff, artifacts, algorithms, personal roles...

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How functions relate to means and ends



"That switch mutes the television."

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How functions relate to means and ends



"That switch mutes the television." ↓ One can *use* the switch to mute the television.

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"That switch mutes the television."
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One can use the switch to mute the television.
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 - distinct of desirability
 - distinct from theory of practical reasoning

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Initial analysis of means-end relations

• An end is some desirable condition – a proposition.

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Some controversies:

- Ends-in-themselves?
- Objects as means?

PDL syntax

Propositional Dynamic Logic is a logic of actions.

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Basic types: • a set **act** of *actions*,

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Basic types:

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 - sequential composition $\alpha; \beta$
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 - Closed under:
 - boolean connectives,
 - dynamic operators $[\alpha]\varphi$, $\langle \alpha \rangle \varphi$.

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Intuitions:

• $[\alpha]\varphi$: after doing α , φ will hold.

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Intuitions:

- $[\alpha]\varphi$: after doing α , φ will hold.
- $\langle \alpha \rangle \varphi$: after doing α , φ might hold.

PDL semantics



Possible world semantics with transition systems for each action α .

PDL semantics



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 $w \xrightarrow{\alpha} w'$ means:

one can reach w' by doing α in w.

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 $\mathbf{w} \models [\alpha] \varphi \quad iff \quad \forall \ \mathbf{w} \xrightarrow{\alpha} \mathbf{w}' \quad \mathbf{w}' \models \varphi.$

PDL semantics



Possible world semantics with transition systems for each action α .

 $w \xrightarrow{\alpha} w'$ means:

one can reach w' by doing α in w.

$$w \models [\alpha]\varphi \quad iff \ \forall \ w \xrightarrow{\alpha} w' \ . \ w' \models \varphi.$$
$$w \models \langle \alpha \rangle \varphi \quad iff \ \exists \ w \xrightarrow{\alpha} w' \ . \ w' \models \varphi.$$

A means is an action α that can realize one's end $\varphi.$

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Two interpretations:



Weak: α might realize φ .

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Weak: α might realize φ . Strong: α will realize φ .

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Two interpretations:



Weak: α might realize φ . $w \models \langle \alpha \rangle \varphi$



Strong: α will realize φ . $w \models [\alpha] \varphi \land \underline{\langle \alpha \rangle} \top$ α can be done.

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Different means to a common end have different degrees of reliability.

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Different means to a common end have different degrees of reliability.

End: Get 12 points with one dart.



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Three different means:

• Throw for 12.



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Efficacy: The degree of reliability of a means to an end.



Efficacy is a measure of likelihoods.



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PDL includes non-determinism, not probabilities.



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Fix (semantic): use *probabilistic* transition structures.

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 $w \xrightarrow[x]{} w'$ means that doing α in w has probability xof resulting in w'.

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<u>Write</u>: $P(w \xrightarrow{\alpha} w') = x$.



Syntactic fix?

Hughes, Esterline, Kimiaghalam Means-end Relations and a Measure of Efficacy



Syntactic fix?

• Probabilistic Computation Tree Logic (pCTL)?

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Syntactic fix?

- Probabilistic Computation Tree Logic (pCTL)?
 - Index dynamic operators, like $[\alpha]$ $\langle \alpha \rangle$

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like
$$[\alpha]_{\geq x}$$
, $\langle \alpha \rangle_{\geq x}$.



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- Probabilistic Computation Tree Logic (pCTL)?
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 - Nesting requires picking x's.

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• Probabilistic PDL?



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- Probabilistic PDL?
 - Truth functional.



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 - Assigns values in [0, 1] to world-formula pairs.

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• Logic in loose sense.



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• Logic in loose sense.

• Fuzzy PDL.

Slogan: Probabilities and fuzziness are different.

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But one can use probabilities to define fuzzy predicates.

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But one can use probabilities to define fuzzy predicates.

Hajek, et al., uses distributions on propositional formulas to define "Probably φ ".

Truth degree of "Probably φ " = $P(\varphi)$.

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$$\llbracket \langle \alpha \rangle \varphi \rrbracket(w) = \sum_{w' \in \mathcal{W}} P(w \xrightarrow{\alpha} w') \cdot \llbracket \varphi \rrbracket(w').$$

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• Like decision theory, we use means for expected outcomes.

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• Like decision theory, we use means for expected outcomes.

• Unlike decision theory, there are no utilities involved.

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• Like decision theory, we use means for expected outcomes.

- Unlike decision theory, there are no utilities involved.
- Elegant treatment of complex ends, like $\langle \alpha \rangle \varphi \wedge \langle \beta \rangle \psi$.

Fuzzy ends An accidental advantage

Weapons are for causing harm.

Means-end relations Efficacy via fuzzy logic Reliability as a fuzzy operator The resulting fuzzy logic

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Weapons are for causing harm. Examples: slingshot, nuke

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This end is fuzzy.



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Fuzzy ends An accidental advantage



Weapons are for causing harm. Examples: slingshot, nuke This end is fuzzy. Fuzzy PDL allows for fuzzy ends. A nuke is more effective in causing harm than a slingshot. (Duh.)

Suppose J and L are cooperative but incommunicado.

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Reliability as a fuzzy operator The resulting fuzzy logic

Extending the logic to other connectives

Suppose J and L are cooperative but incommunicado.

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He wants to ensure that *L* will succeed, whichever she chooses. *End*: $\langle m \rangle P \land \langle n \rangle Q$. *Aim*: maximize min{ $[\langle m \rangle P](w), [\langle n \rangle Q](w)$ }.

The resulting fuzzy logic

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He wants to ensure that L will succeed, whichever she chooses. End: $\langle m \rangle P \land \langle n \rangle Q$. Aim: maximize min{ $[\![\langle m \rangle P]\!](w), [\![\langle n \rangle Q]\!](w)$ }.

 $\llbracket \varphi \land \psi \rrbracket(w) = \min \{\llbracket \varphi \rrbracket(w), \llbracket \psi \rrbracket(w) \}$

On formulas

$$\begin{split} \llbracket \langle \alpha \rangle \varphi \rrbracket(w) &= \sum_{w' \in \mathcal{W}} P(w \xrightarrow{\alpha} w') \cdot \llbracket \varphi \rrbracket(w') \\ \llbracket \varphi \land \psi \rrbracket(w) &= \min \{ \llbracket \varphi \rrbracket(w), \llbracket \psi \rrbracket(w) \} \\ \llbracket \varphi \lor \psi \rrbracket(w) &= \max \{ \llbracket \varphi \rrbracket(w), \llbracket \psi \rrbracket(w) \} \\ \llbracket \neg \varphi \rrbracket(w) &= 1 - \llbracket \varphi \rrbracket(w) \\ \llbracket \varphi \to \psi \rrbracket(w) &= \begin{cases} 1 & \text{if } \llbracket \varphi \rrbracket(w) \le \llbracket \psi \rrbracket(w), \\ \llbracket \psi \rrbracket(w) &= \text{else;} \end{cases} \end{split}$$

On formulas

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On actions

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The semantics of fuzzy PDL

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 $\left[\begin{array}{c} [\varphi \cup \psi](w)(w') \\ [\varphi^*](w)(w') \end{array} \right\} \text{ undefined.}$

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 - Axioms:
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Logical properties Completeness

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Logical properties Completeness

I wish.

Hughes, Esterline, Kimiaghalam Means-end Relations and a Measure of Efficacy

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Logical properties Completeness

I wish.

But not with these semantics.

Ongoing work...

• Include non-deterministic features (in paper).

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Thank you.

Concerns:

• Primary: Adding probabilities to transitions.

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<u>Aims</u>:

• Keep PDL as language for means-end relations.

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Proposal: Interpret PDL as fuzzy logic.

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